

Blind Detection of Polar Codes

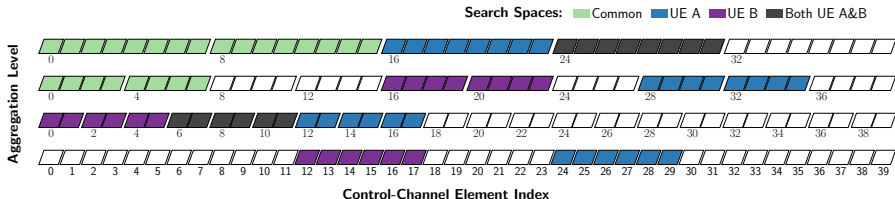
Pascal Giard, Alexios Balatsoukas-Stimming, and Andreas Burg

Telecommunications Circuits Laboratory, EPFL



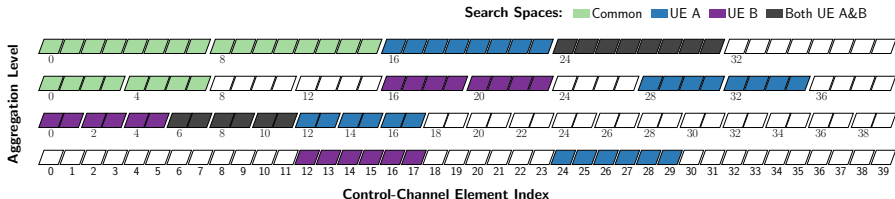
Control Channel Detection for 5G

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- Don't know in advance...
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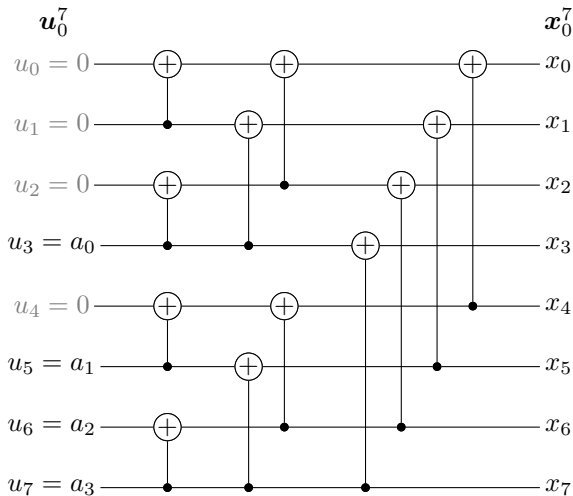
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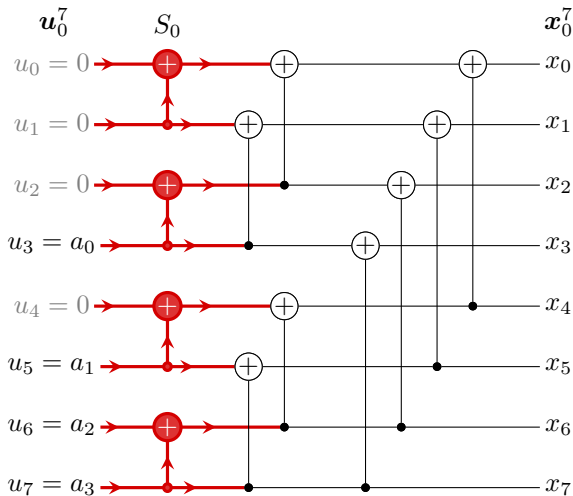
How should we define that metric?

About Polar Codes

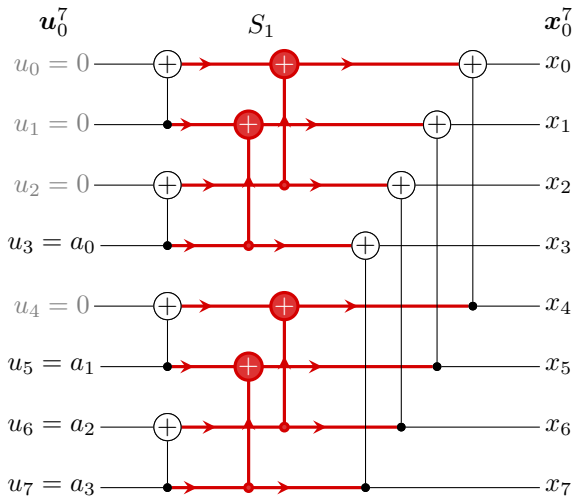
Polar Codes – Encoding



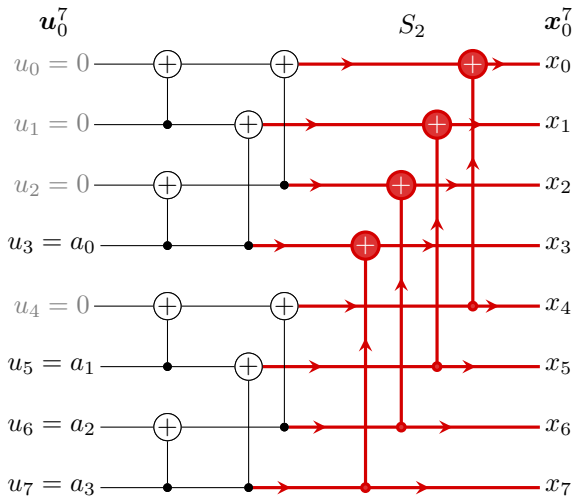
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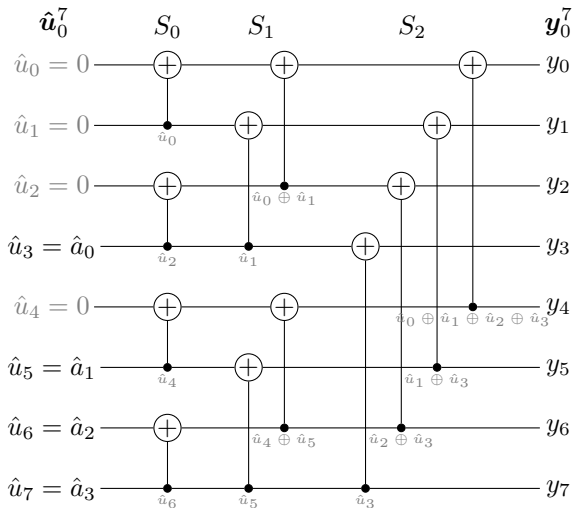
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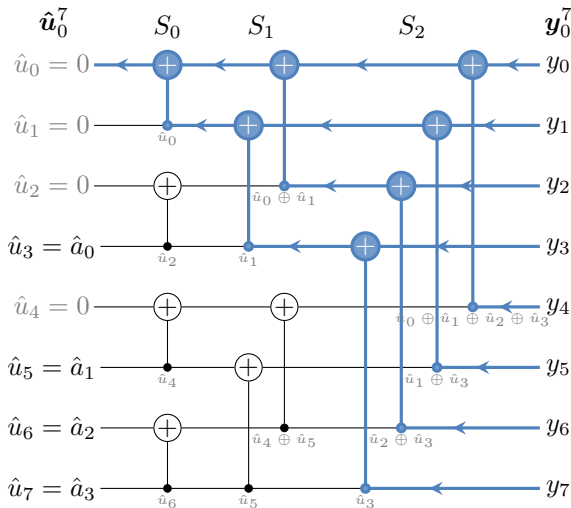
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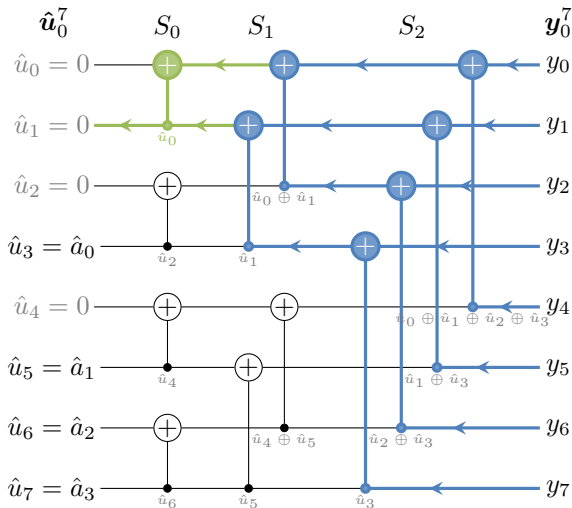
Polar Codes – Successive-Cancellation Decoding



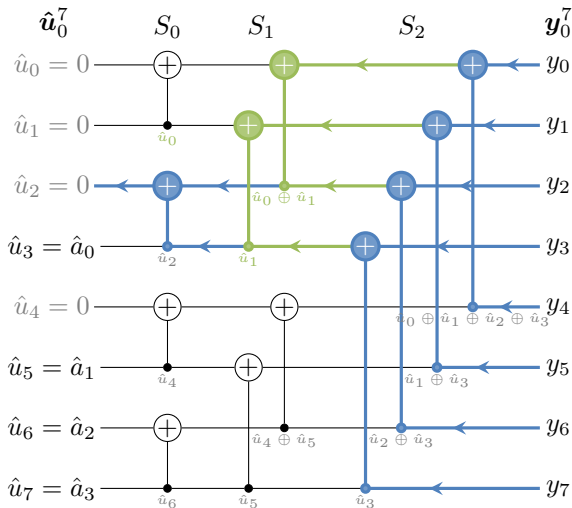
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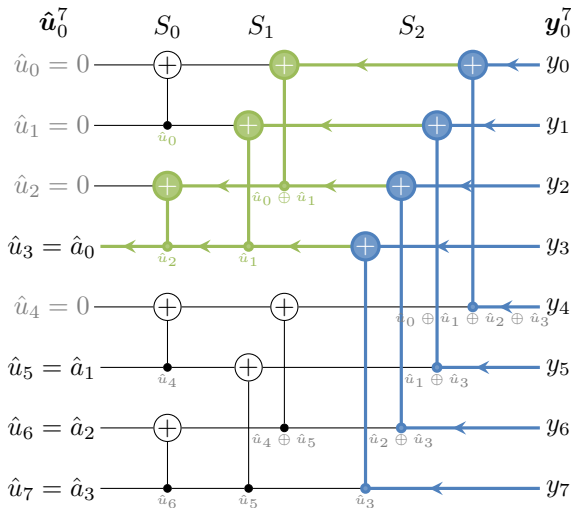
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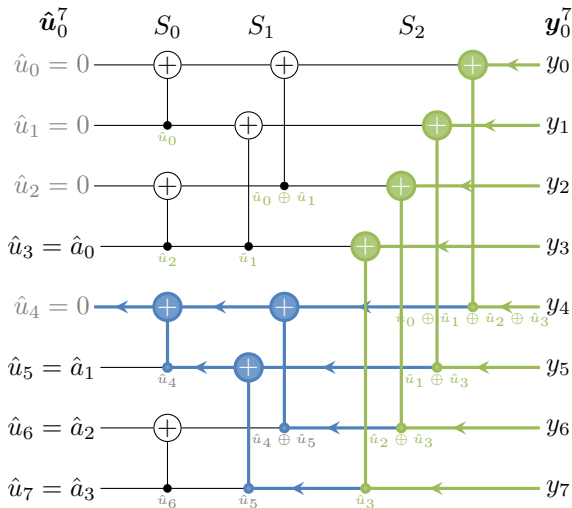
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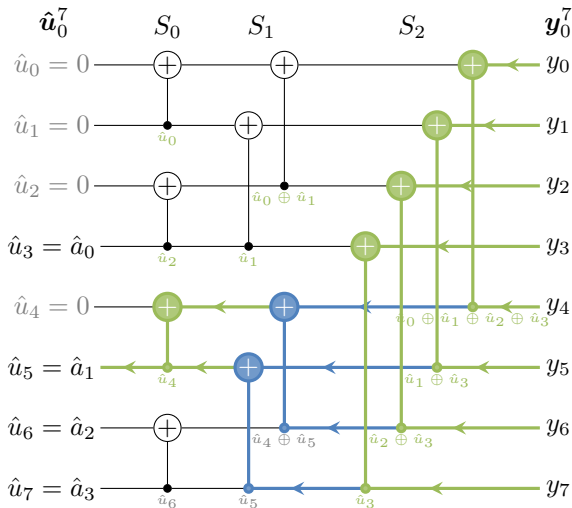
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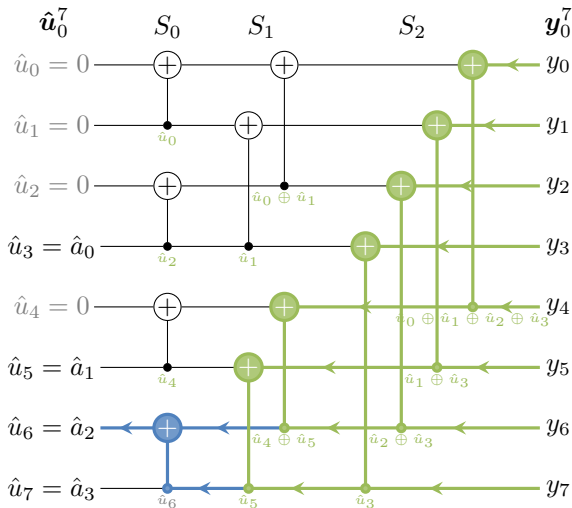
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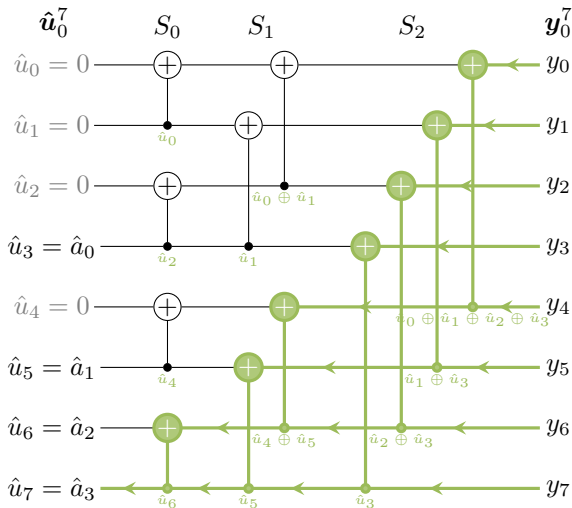
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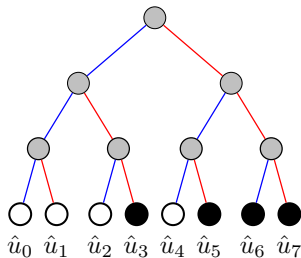
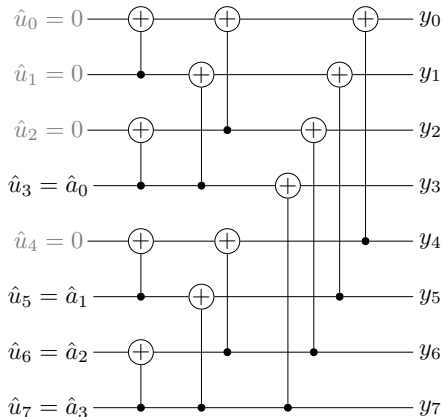
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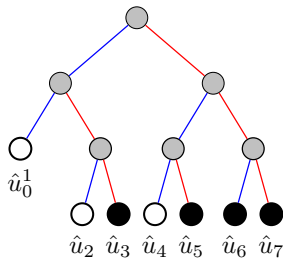
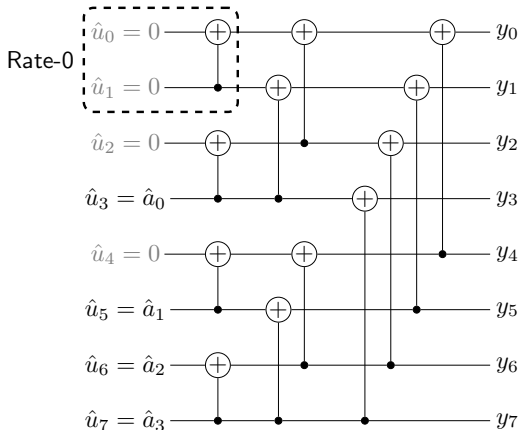
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Polar Code as a Concatenation of Codes

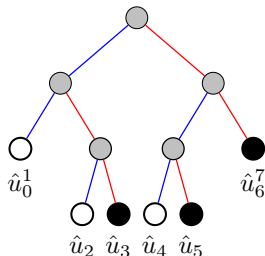
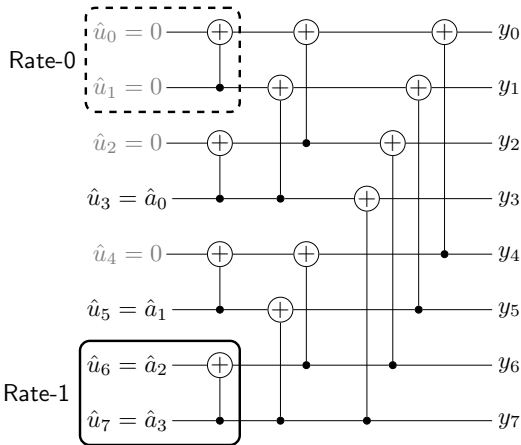


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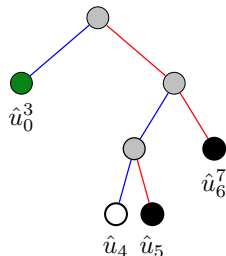
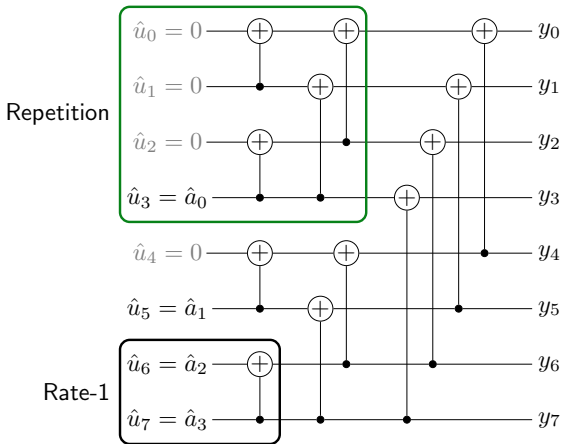
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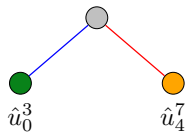
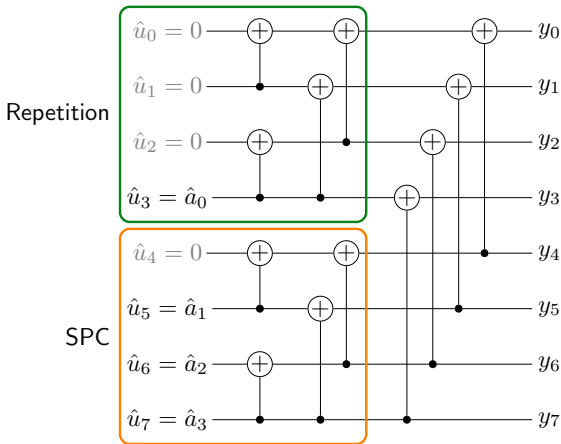
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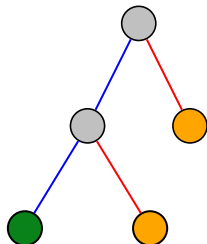
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Proposed Blind-Detection Method

Detection Metric

- Exploit the inherent structure of three constituent-code types to compute a detection metric
- The bigger the value of the detection metric, the more likely a received block is encoded with the expected polar code
- Use the detection metric to determine which blocks are forwarded to the next stage, i.e., the more complex decoder



Proposed Steps

- 1 Conduct (complete or partial) Fast-SSC decoding on all candidates

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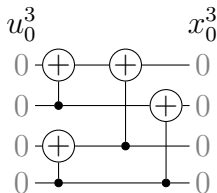
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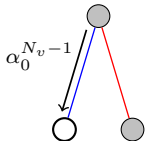
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- 3 Pass on a subset of the best candidates to a more complex SCL decoder

Update Rules – Rate-0 “Code”

- Estimated bit vector \hat{u} known *a priori* to be solely made of frozen bits, i.e., to be an all-zero vector



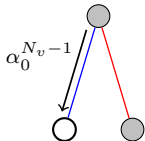
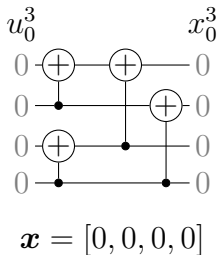
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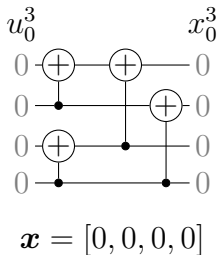
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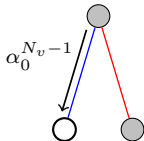
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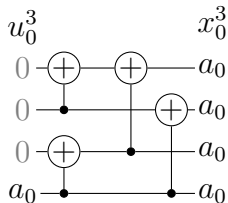
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- Scaling factor to normalize w.r.t. to constituent code length

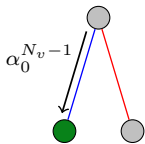


Update Rules – Repetition Code

- The only information bit is repeated on all outputs

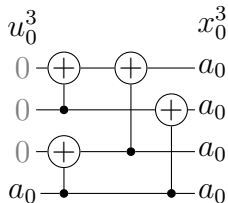


$$\mathbf{x} = [a_0, a_0, a_0, a_0]$$

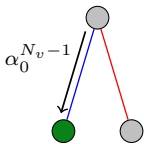


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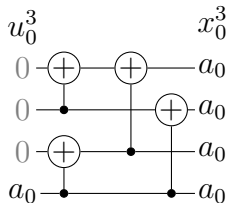
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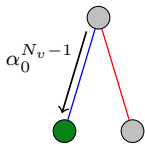
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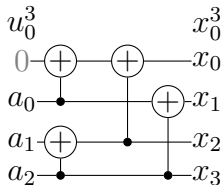
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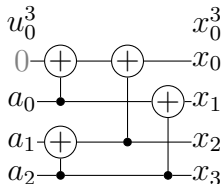


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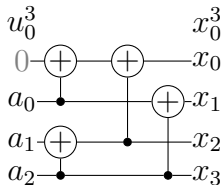


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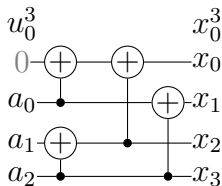


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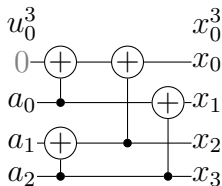
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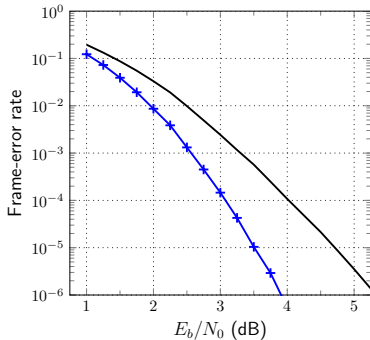
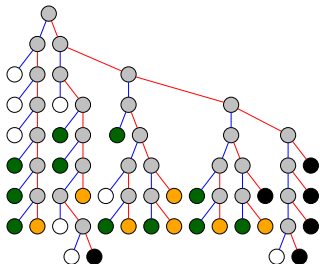
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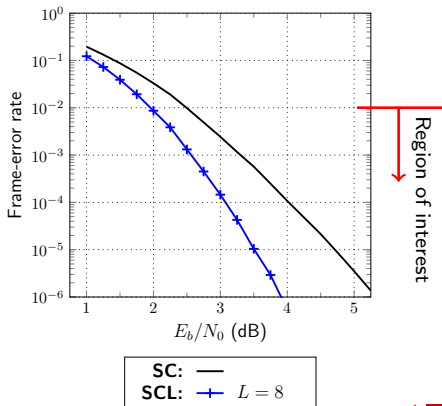
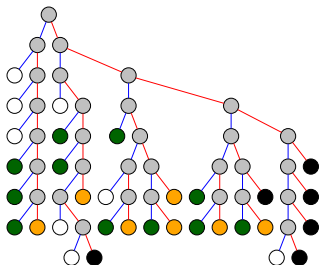
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SC: —
SCL: + $L = 8$

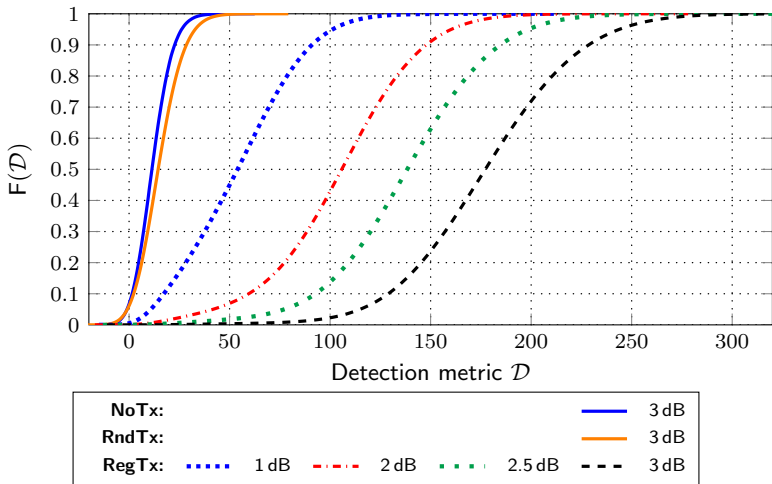
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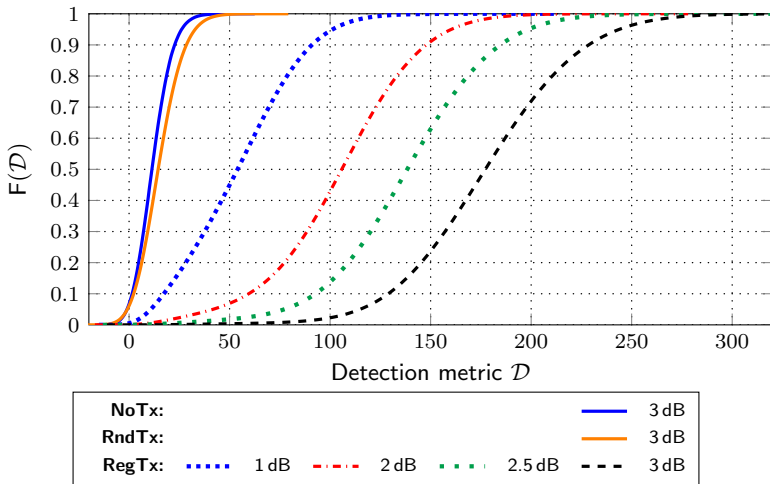
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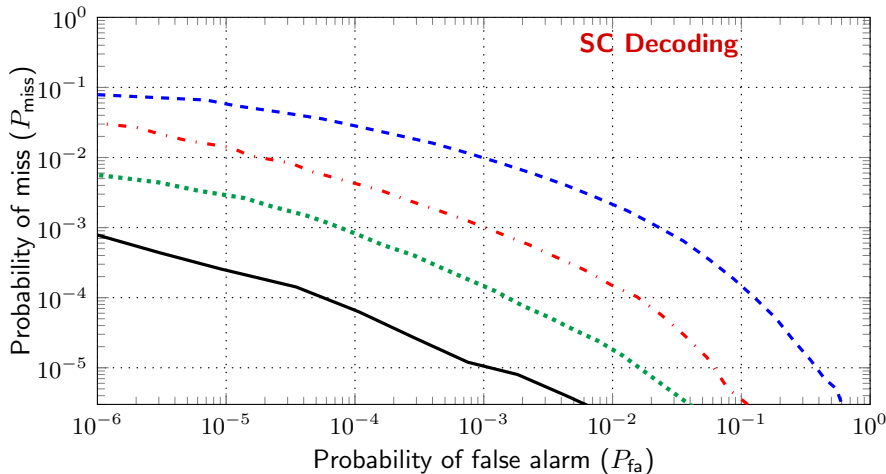
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Metric starts to be really good at 2 dB, the region of interest!

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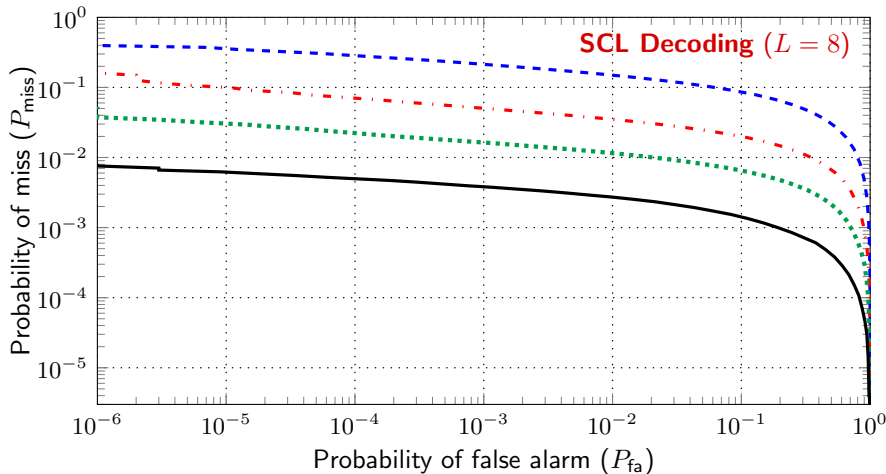
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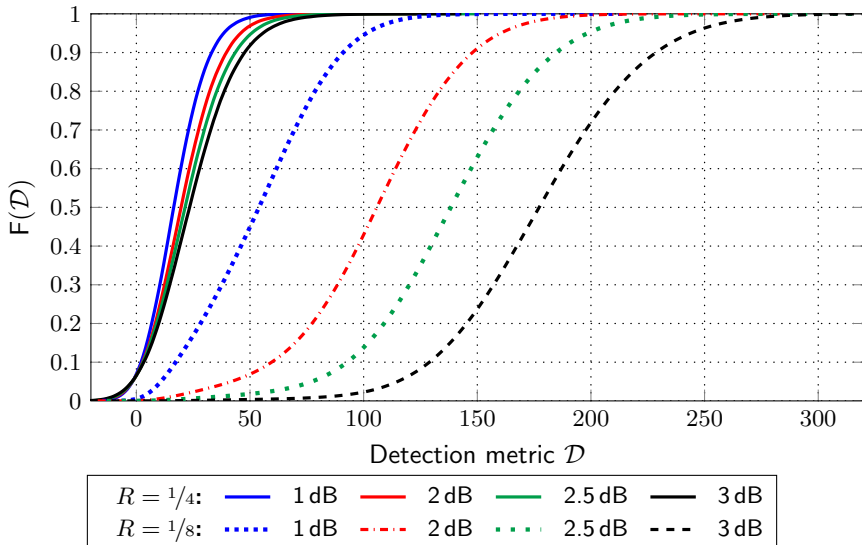
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 - Built from a fast-SSC decoder
 - Exploits the structure of constituent codes
- Allows to quickly reduce a list of candidates to a tractable number

Thank you for listening!

Bonus Slides

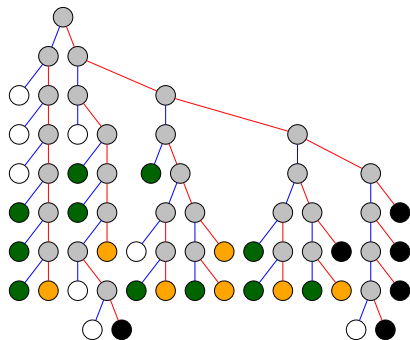
Correlated Input – Different Rate

Experimental CDFs of \mathcal{D} ; detecting for $R = 1/8$

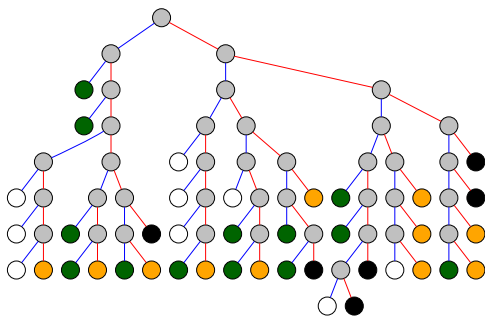


Different Rate – Different Decoder Tree

$$R = 1/8$$



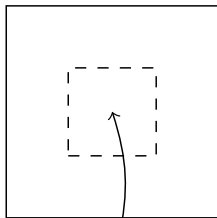
$$R = 1/4$$



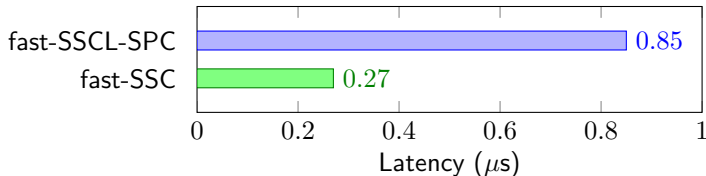
Complexity of the Detector

- **Worst case** time and area complexities approach that of a fast-SSC decoder
- May not be necessary to run the detector on the complete decoder tree
- Area complexity of a fast-SSC decoder is much lower than that of an SCL decoder for $L = 8$

fast-SSCL-SPC



fast-SSC



Hashemi, Condo, and Gross., "Fast and Flexible Successive-Cancellation List Decoders for Polar Codes," IEEE TSP, 2017.

Giard, Balatsoukas-Stimming, Sarkis, Thibault, and Gross., "Fast Low-Complexity Decoders for Low-Rate Polar Codes," Springer JSPS, 2016.